

draw Backoff slots $\rightarrow x \in [0, C], x \in \mathcal{N}$

1) a) $\Pr[X=n] = \frac{1}{C+1}$

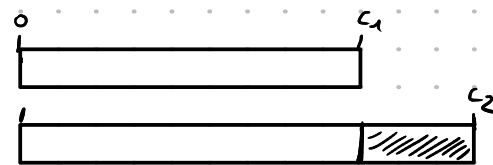
b) $\Pr[X \leq n] = \sum_{i=0}^n \Pr[X=i] = \frac{n+1}{C+1}$

c) Nodes 1, 2 where $C_1 = C_2 = C$
 $\left. \begin{array}{l} \uparrow \text{ } y \text{ slots} \\ x \text{ slots} \end{array} \right\} \text{Collision iff } y = x$

$$\begin{aligned} \Pr[y=x] &= \sum_{i=0}^C \Pr[X_1=i \wedge X_2=i] \\ &= \sum_{i=0}^C \Pr[X_1=i] \Pr[X_2=i] \\ &= \sum_{i=0}^C \frac{1}{C+1} \cdot \frac{1}{C+1} = \sum_{i=0}^C \left(\frac{1}{C+1}\right)^2 \\ &= (C+1) \cdot \left(\frac{1}{C+1}\right)^2 = \frac{1}{C+1} \end{aligned}$$

d) Nodes 1, 2 choose x, y slots. $x \in [0, C_1] \mid C_1 \neq C_2$
 $y \in [0, C_2]$

$$\Pr[\text{"collision"}] = \sum_{i=0}^{\min(C_1, C_2)} \Pr[X_1=i \wedge X_2=i]$$



let without loss of generality $C_2 \geq C_1$

$$\begin{aligned} &\Rightarrow \sum_{i=0}^{C_1} \Pr[X_1=i] \cdot \Pr[X_2=i] \\ &= \sum_{i=0}^{C_1} \frac{1}{C_1+1} \cdot \frac{1}{C_2+1} = (C_1+1) \cdot \frac{1}{(C_1+1)(C_2+1)} \\ &= \frac{1}{C_2+1} \end{aligned}$$

e) Node 1, 2 $\left\{ \begin{array}{l} x \in [0, c_1] \\ y \in [0, c_2] \end{array} \right. \Bigg| c_1 \leq c_2$

$$c_1 \leq c_2$$

should be $c_1 < c_2$, otherwise case $\Pr[x=c_1] \Pr[y>c_1]$ broken

Node 1 wins iff $x < y$

$$\Pr[x < y] = \sum_{i=0}^{c_1} \Pr[x=i] \Pr[y > i]$$

$\hookrightarrow 1 - \Pr[y \leq i]$

$$= \sum_{i=0}^{c_1} \frac{1}{c_1+1} \left(1 - \frac{i+1}{c_2+1} \right)$$

$$= \sum_{i=0}^{c_1} \frac{1}{c_1+1} - \frac{i+1}{(c_1+1)(c_2+1)}$$

$$= \sum_{i=0}^{c_1} \frac{1 \cdot \cancel{(c_2+1)}}{(c_1+1) \cancel{(c_2+1)}} - \frac{i+1}{(c_1+1)(c_2+1)}$$

$$= \sum_{i=0}^{c_1} \frac{c_2}{(c_1+1)(c_2+1)} - \frac{i}{(c_1+1)(c_2+1)}$$

$$= \frac{\cancel{(c_1+1)} \cdot c_2}{\cancel{(c_1+1)}(c_2+1)} - \sum_{i=0}^{c_1} \frac{i}{(c_1+1)(c_2+1)} \Bigg| \sum_{i=0}^{c_1} i = \sum_{i=1}^{c_1} i$$

$$= \frac{c_2}{c_2+1} - \frac{\cancel{c_1} \cdot \cancel{(c_1+1)}}{2 \cancel{(c_1+1)}(c_2+1)}$$

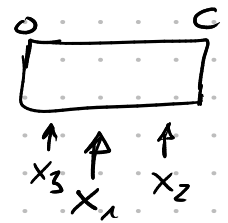
$$= \frac{2(c_2)}{2(c_2+1)} - \frac{c_1}{2(c_2+1)} = \frac{2c_2 - c_1}{2(c_2+1)}$$

f) expected duration of cont. phase T_n for n nodes
 where $C_i = C_j, \forall i, j \in X$ $X = \{0, 1, \dots, C\}$

$$E[X] = \sum_{k=0}^{\infty} k \Pr[X=k] = \sum_{k=0}^{\infty} 1 - \Pr[X \leq k] = \boxed{\sum_{k=0}^{\infty} \Pr[X > k]}$$

$$E[T_n] = E[\min\{X_i \mid i \in \{1, \dots, n\}\}]$$

$$= \sum_{k=0}^C \Pr[X_i > k \mid i \in \{1, \dots, n\}]$$



$$= \sum_{k=0}^C \prod_{j=1}^n \Pr[X_j > k] = \sum_{k=0}^C \prod_{j=1}^n 1 - \Pr[X \leq k]$$

$$= \sum_{k=0}^C \prod_{j=1}^n 1 - \frac{k+1}{C+1} = \sum_{k=0}^C \prod_{j=1}^n \frac{1(C+1) - (k+1)}{C+1} = \sum_{k=0}^C \prod_{j=1}^n \frac{C-k}{C+1}$$

← swap sum direction

$$= \sum_{k=0}^C \prod_{j=1}^n \frac{C-k}{C+1} = \sum_{k=0}^C \left(\frac{C-k}{C+1}\right)^n = \sum_{k=0}^C \left(\frac{k}{C+1}\right)^n$$

g) exp. dur. of cont. phase T for 2 nodes with $C_1 \neq C_2$

$$E[T] = E[\min\{X_1, X_2\}]$$

$$= \sum_{k=0}^{\min\{C_1, C_2\}} \Pr[X_1 > k \wedge X_2 > k]$$

$$= \sum_{k=0}^{\min\{C_1, C_2\}} \Pr[X_1 > k] \Pr[X_2 > k]$$

$$= \sum_{k=0}^{\min\{C_1, C_2\}} (1 - \Pr[X_1 \leq k]) (1 - \Pr[X_2 \leq k])$$

$$= \sum_{k=0}^{\min\{C_1, C_2\}} \left(1 - \frac{k+1}{C_1+1}\right) \left(1 - \frac{k+1}{C_2+1}\right)$$

$$= \sum_{k=0}^{\min\{C_1, C_2\}} \left(\frac{C_1+1}{C_1+1} - \frac{k+1}{C_1+1}\right) \left(\frac{C_2+1}{C_2+1} - \frac{k+1}{C_2+1}\right)$$

$$= \sum_{k=0}^{\min\{C_1, C_2\}} \frac{C_1-k}{C_1+1} \cdot \frac{C_2-k}{C_2+1} = \dots$$