## Network Coding (NC)

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#### What is [libmoep](#page-2-0)?

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 $1$ ibmoep<sup>1</sup> is a library written in C that allows to inject

- cooked IEEE 802.11 frames (native mode),
- frames based on a proprietary, extensible frame format (moep 802.11) to develop and evaluate custom link-layer protocols and
- various other frametypes, and
- it supports various interfaces (WLAN, Ethernet, TAP, Unix sockets).

It has been primarily developed by Maurice Leclaire, a former staff member of the chair.

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- various other frametypes, and
- it supports various interfaces (WLAN, Ethernet, TAP, Unix sockets).

Why not opening raw sockets? ...1ibmoep uses raw sockets but:

- it hides most of the complexity of
	- creating monitor mode interfaces.
	- setting interface parameters,
	- parsing radiotap headers, etc.
- and allows a convenient way to *pair* a monitor interface with a TAP interface.

It has been primarily developed by Maurice Leclaire, a former staff member of the chair.

### What is [libmoep](#page-2-0)?

#### Example: ptm

The ptm (PTM stands for packet transfer module) is

- the most simple kind of module using libmoep to
- relay packets by
	- accepting IEEE 802.3 frames over a virtual Ethernet interface (tap0),
	- converting those frames to a custom format suitable for wireless transmission,
	- sending those frames over a monitor interface and
	- translating incoming frames from the monitor interface back to valid IEEE 802.3 frames.



# пm

## What is [libmoep](#page-2-0)?

#### Example: ptm

The TAP interface presents itself like a physical Ethernet device, i.e.,

- it has a MAC address and
- can be assigned IP(v6) addresses:



## What is [libmoep](#page-2-0)?

#### Example: ptm

The TAP interface presents itself like a physical Ethernet device, i.e.,

- it has a MAC address and
- can be assigned IP(v6) addresses:

tap0: <BROADCAST.UP.LOWER\_UP> mtu 1500 qdisc pfifo\_fast ... 2 link / ether 06:36:10:3 e : a8 : b0 brd ff : ff : ff : ff : ff : ff 3 inet 10.0.0.1/24 brd 10.0.0.255 scope global tap0 valid lft forever preferred lft forever 5 inet6 fe80 ::436:10 ff : fe3e : a8b0 /64 scope link

Advantages:

- Applications are completely unaware of the translation.
- It works with any kind of traffic (even ARP).
- We have any control about the radio interface we can ever have without writing custom device drivers.

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#### [How is it implemented?](#page-9-0)

1. Create tap and monitor devices:

```
if (! (tap = moep dev ieee8023 tap open (args addr , & args ip , 24
2 args.mtu + sizeof(struct ether header )))) {
3 fprintf ( stderr , " ptm : error : %s\ n " , strerror ( errno ));
          return -1;
5 }
6
7 if (!(rad = moep_dev_moep80211_open (args.rad, args.freq,
8 MOEP80211 CHAN WIDTH 20 .
9 0 0 args mtu + radiotap len (-1) +
10 sizeof ( struct moep80211_hdr ) +
11 sizeof ( struct moep_hdr_pctrl )))) {
12 fprintf( stderr, " ptm: error: % s\n", strerror(errno));
13 moep_dev_close (tap);
14 return -1;
15 }
```
2. Set rx\_handler for both devices that will be used as callbacks upon frame arrival:

```
moep dev set rx handler ( tap , taph ) :
```
moep\_dev\_set\_rx\_handler(rad, radh);

3. *Pair* both devices and turn control to libmoep:

```
moep_dev_pair (tap, rad);
```

```
moep_run( sigh , NULL ) ;
```
### [How is it implemented?](#page-9-0)

- The call to moep\_run() turns control to libmoep.
- The internal event loop is essentially a wrapper for epoll.
- Depending on which interface a frame is received, the appropriate handler is called:
	- If a frame arrives at the TAP interface, taph() is called and the received frame is passed to this handler.
	- The handler can translate the frame to a suitable format and schedule it for transmission on the radio interface.

#### [How is it implemented?](#page-9-0)

- The call to moep run() turns control to libmoep.
- The internal event loop is essentially a wrapper for epoll.
- Depending on which interface a frame is received, the appropriate handler is called:
	- If a frame arrives at the TAP interface, taph() is called and the received frame is passed to this handler.
	- The handler can translate the frame to a suitable format and schedule it for transmission on the radio interface.

#### **Do we have to turn complete control over to libmoep?**

- Of course not.
- There is moep wait(), which works just like epoll wait() but still supports rx handlers.
- You can even configure libmoep to use another custom epoll wait() compatible function internally.

However, you will probably never need that.

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#### [moep 80211 frame format](#page-13-0)

There are two different ways to create radio interfaces:

- moep\_dev\_ieee80211\_open()
	- Frames passed to the rx\_handler will be ordinary IEEE 802.11 frames, including their link-layer headers.
	- The radiotap header will be a moep80211 radiotap since ieee80211 radiotap sucks.<sup>1</sup>
- moep dev moep80211 open()
	- Frames passed to the rx\_handler will be in a custom format that is based on the generic IEEE 802.11 header for data frames.
	- The radiotap header is again moep80211\_radiotap.

There are more functions to open / create TAP and Ethernet interfaces (and even Unix sockets).

In all cases, a frame is represented by the type moep\_frame\_t, which is a typedef of a pointer to a struct moep\_frame.

- The members of this struct are an implementation detail and not accessible.
- You cannot modify a moep frame t directly, use the interface functions.

The difference between moep80211\_radiotap and ieee80211\_radiotap is basicalle that the former one is fully expanded, i.e., all options allocate memory even if the present mask does not specify them.

#### [moep 80211 frame format](#page-13-0)

Getting the headers of a moep frame t:

1 // Returns the radiotap header

struct moep80211 radiotap \* moep frame\_radiotap ( moep frame t frame );

- 3 // Returns the IEEE80211 header ( generic format , your have to parse it)
- 4 struct ieee80211 hdr gen \* moep frame\_ieee80211\_hdr ( moep frame t frame );
- 5 // Returns the moep80211\_hdr common to all our custom frames
- 6 struct moep80211 hdr \* moep frame moep80211 hdr ( moep frame t frame );
- 7 // Returns the IEEE802 .3 header (in case of Ethernet frames )
- 8 struct ether header \* moep\_frame\_ieee8023\_hdr ( moep\_frame\_t frame );

Transmit a frame:

 $int$  moep dev  $tx$  (moep dev  $t$  dev, moep frame  $t$  frame);

If you want to convert a moep frame t to a buffer or create a moep frame t from a buffer, you can use the following functions:

int moep frame encode ( moep frame t frame, u8 \*\* buf, size t buflen );

int moep frame decode ( moep frame t frame, u8 \* buf, size t buflen );

If you want to transmit a manually created frame from a buffer, you may also use:

int moep dev tx raw (moep dev t dev, u8 \* buf, size t buflen);

```
1 struct moep80211_hdr {
2 u16 frame_control:
3 u16 duration_id ;
4 u8 ra [ IEEE80211_ALEN ];
5 u8 ta [ IEEE80211_ALEN ];
6 u32 disc ;
7 u16 txseq ;
8 u16 seq_ctrl; /* unused/reserverd */
9 } __attribute__((packed));
```

```
1 struct moep80211 hdr {
2 u16 frame_control:
3 u16 duration_id ;
4 u8 ra [ IFFF80 211 ALEN ]:
5 u8 ta [ IEEE80211_ALEN ];
6 u32 disc ;
7 u16 txseq ;
8 u16 seq ctrl: /* unused/reserverd */
9 } __attribute__((packed));
```
- frame\_control has the same meaning as for ordinary IEEE 802.11 frames.
- We set it to FTYPE\_DATA | STYPE\_DATA for **all** of our frames to avoid unexpected behavior of hardware.

```
1 struct moep80211_hdr {
2 u16 frame_control:
3 u16 duration_id ;
4 u8 ra [ IEEE80211_ALEN ];
5 u8 ta [ IEEE80211_ALEN ];
6 u32 disc ;
7 u16 txseq ;
8 u16 seq ctrl: /* unused/reserverd */
9 } __attribute__((packed));
```
- duration\_id may be interpreted by other STAs.
- We set it to zero for now.

```
1 struct moep80211 hdr {
2 u16 frame_control:
3 u16 duration_id ;
4 u8 ra [ IFFF80 211 ALEN ]:
5 u8 ta [ IEEE80211_ALEN ];
6 u32 disc ;
7 u16 txseq ;
8 u16 seq ctrl: /* unused/reserverd */
9 } __attribute__((packed));
```
- ra is the 6 B receiver address of this frame.
- If we exploit the wireless broadcast advantage, we set it to the MAC broadcast address.

```
1 struct moep80211 hdr {
2 u16 frame_control:
3 u16 duration_id ;
4 u8 ra [ IFFF80 211 ALEN ]:
5 u8 ta [ IEEE80211_ALEN ];
6 u32 disc ;
7 u16 txseq ;
8 u16 seq ctrl: /* unused/reserverd */
9 } __attribute__((packed));
```
- ta is the 6 B transmitter address of this frame.
- This is **not** the MAC of our wireless interface but of the tap interface. Think about it!

```
1 struct moep80211 hdr {
2 u16 frame_control:
3 u16 duration_id ;
4 u8 ra [ IFFF80 211 ALEN ]:
5 u8 ta [ IEEE80211_ALEN ];
6 u32 disc ;
7 u16 txseq ;
8 u16 seq ctrl: /* unused/reserverd */
9 } attribute ((packed));
```
- disc is a 4 B field that we call frame discriminator.
- In IEEE 802.11 data frames this would be the third MAC address.
- We choose a value that should be invalid as MAC address.
- This way we can differentiate our own frames from normal IEEE 802.11 traffic.

```
1 struct moep80211 hdr {
2 u16 frame_control:
3 u16 duration_id ;
4 u8 ra [ IFFF80 211 ALEN ]:
5 u8 ta [ IEEE80211_ALEN ];
6 u32 disc ;
7 u16 txseq ;
8 u16 seq ctrl: /* unused/reserverd */
9 } __attribute__((packed));
```
- txseq are the latter 2 B of the third MAC address in IEEE 802.11 data frames.
- We use it as per-node TX sequence number, e.g. to estimate erasure probabilities.

```
1 struct moep80211 hdr {
2 u16 frame_control:
3 u16 duration_id ;
4 u8 ra [ IFFF80 211 ALEN ]:
5 u8 ta [ IEEE80211_ALEN ];
6 u32 disc ;
7 u16 txseq ;
8 u16 seq ctrl: /* unused/reserverd */
9 } attribute ((packed));
```
- seq\_ctr1 is fragment number / sequence number of the normal IEEE 802.11 data frame header.
- Problem with this field is that the NIC's driver may play with it.
- It is safer to set it to zero and to ignore it on reception.

#### [moep 80211 frame format](#page-13-0)

#### Extension headers

We use extension headers to resemble different frame types, e.g. the packet control header:

- After the moep hdr at **least one** extension header must follow.
- Bit 7 in the extension header's type field indicates whether another extension header follows.
- Type and length field precisely specify the extension header, and allow anyone to skip unknown extension headers.



```
struct moep hdr_pctrl {
2 struct moep hdr ext hdr:
3 u16 type: // corresponding to the Ethertype
4 u16 len ; // explicit length of the frame 's payload
5 } attribute ((packed));
6
7 struct moep_hdr_ext {
8 u8 type; // type of the extension header, e.g. MOEP_HDR_PCTRL
9 u8 len: // total length of the extension header
10 } attribute (( packed )):
```
- Extension headers are part of the l2\_header in the private struct moep\_frame.
- How exactly extension headers are stored within a typedeffed moep\_frame\_t is not your business.
- Extension headers are part of the 12\_header in the private struct moep\_frame.
- How exactly extension headers are stored within a typedeffed moep\_frame\_t is not your business.

Just let libmoep do it for you:

- moep frame add moep hdr ext() Add a new extension header to an existing frame.
- moep\_frame\_set\_moep\_hdr\_ext() Replace an existing extension header by a new one
- moep frame del moep hdr ext() Delete an extension header.
- moep frame moep hdr ext()

Get a pointer to a specific extension header (or NULL if it does not exist).

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#### [Blocking interfaces](#page-27-0)

As handlers must never perform blocking operations, we have to deal with the possibility that

- we read a frame from the TAP interface (taph()) but
- cannot write it to the monitor interface.

(The reverse way may occur in theory but is of no practical interest.)

#### **Possible solutions:**

- 1. Discard the packet
- 2. Buffer the packet
- 3. Do not read from the TAP interface in the first place, eventually leading to frame drops in the kernel

If we use random linear network coding, we have buffers and thus

- read from the TAP interface until our internal buffers are filled and
- block the TAP interface until we have free space again.

libmoep offers the possibility to "block" and "unblock" an interface with moep\_dev\_set\_rx\_status().

Normally, an interface should be "blocked" (receiving) in case another interface is blocking transmissions. This event is signalled by libmoep through a callback handler, which can be installed with moep dev set tx status callback().

In the ptm example this combination is automatically installed with moep dev pair(). In more complex scenarios involving buffers, we have to write our own callbacks handling the buffer status.

**Note:** This does not guarantee that no frames are dropped.

- The TAP interface buffers frames on its own, and will eventually drop frames under load.
- Linux does not guarantee that frames scheduled for transmission are indeed transmitted.

**However:** Once a frame is read using libmoep, we can guarantee that it reaches its destination. But we have to ensure it ourselves.

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### [Handling other file descriptors](#page-30-0)

#### Example: tranmit beacon frames in regular time intervals

```
if ((bcn_timer = timerfd_create(ClOCK_MONOTONIC, TED_NONBLOCK)) < 0) {
             fprintf ( stderr , " ptm : error : % s\n" , strerror ( errno ) ) :
            return -1:
 4 }
5
6 interval it interval tv sec = args heacon / 1000;
7 interval . it interval . tv_nsec = ( args . beacon % 1000) * 1000000:
8 interval it value tv_sec = interval it_interval tv_sec ;
9 interval it value tv_nsec = interval it_interval tv_nsec ;
10 if ( timer fd set time ( bcn timer , 0 , & interval , NULL ) ) {
11 fprintf ( stderr , " ptm : error : %s\ n " , strerror ( errno ));
12 close ( bcn_timer );
13 return -1;
14 }
15
16 if (!(bcn_callback = moep_callback_create (bcn_timer, send_beacon, NULL, EPOLLIN))) {
17 fprintf ( stderr, " ptm: error: % s\n", strerror ( errno ));
18 close (bcn_timer):
19 return -1;
20
```
- Create a timer\_fd (file descriptor) bound to a monotonic clock (nonsettable monotonically increasing clock that measures time from some unspecified point, not affected by leap seconds)
- Prepare a struct timespec (two 64 bit integers) representing the interval
- Set the timer's interval
- Register the the timer and its callback in libmoep

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#### [libmoepcommon](#page-32-0)

libmoepcommon is not part of libmoep, but a separate header-only library for common tasks:

- Modified version of the Linux kernel's list implementation
- Implementation of timeouts using event file descriptors
- Macros to calculate sum / difference between struct timespec instances
- Various small functions, e. g. to compare / test MAC addresses, hexdumps, logging, common math operations, etc.

It is part of the network coding module (NCM):

- 1 . 2 | - - benchmark . h
- $1 -$  list  $h$
- 4 | - timeout . h
- $1 -$  types  $h$
- $6$   $|--$  util
- 7 | | | -- alignment.h
- 8 | | | -- assertion .h
- $9$  |  $1 -$  hexdump.h
- $10$   $1 1$   $10$ g, h
- $11 \t -1 \t -2$  mac. h
- 12 | | - maths . h
- 13 | | | -- timespec.h
- $14$   $1- util.h$

#### <span id="page-34-0"></span>[The Linux kernel's list](#page-34-0)

list.h contains a (slightly modified) version of the Linux kernel's list.

• Your list elements are structs:

```
1 struct neighbor {
2 struct list head list:
3 u8 hwaddr [ IEEE80211_ALEN ]:
4 };
```
#### • Create a new list and a new element:

```
1 LIST HEAD ( nblist ):
```

```
2 struct neighbor *nh = calloc(1, sizeof (struct neighbor));
```

```
3 list add (& nb - > list , & nblist):
```
#### • Search for an element in the list:

```
1 struct neighbor *cur:
2 list_for_each_entry ( cur, & nblist, list) {
3 if (0 == memcmp (cur -> hwaddr, hwaddr, IEEE80211_ALEN))
           return cur;
5 }
6 return NULL ;
```
#### [The Linux kernel's list](#page-34-0)

• Remove a given element from the list:

```
listdel (&nb \rightarrow list);
2 free(nb);
```
• Search for an element that shall be removed or remove multiple elements:

```
1 struct neighbor *cur. *tmp:
2 list_for_each_entry_safe ( cur , tmp , & nblist , list ) {
3 if ( is to be removed ( cur )) {
4 list del (& cur - > list);
5 free ( cur );
6 }
7 }
```
**Warning:** removing a list element while iterating with list\_for\_each\_entry invalidates list pointers. Use list\_for\_each\_entry\_safe instead.

#### <span id="page-36-0"></span>**[Timeouts](#page-36-0)**

timeout.h allows you to register a callback that is executed when the timeout times out.

• A timeout is internally represented by:



The callback must be given as function pointer of the following type:

```
typedef int (* timeout cb t)( timeout t, u32, void *);
```
Create a new timeout:

```
if (0 > timeout create ( CLOCK_MONOTONIC , & logt , state_log_cb , NULL ))
2 DIE ("timeout create () failed: %s", strerror ( errno ));
    timeout_settime ( log t , 0, timeout_msec ( LOG_INTERVAL , LOG_INTERVAL ));
```
- To make sure timeouts are signalled, you need to pass a signal handler to moep\_run.
- Inside that handler you have to handle SIGRTMIN.
- timeout\_exec((void \*)siginfo.ssi\_ptr, siginfo.ssi\_overrun);

#### **[Timeouts](#page-36-0)**

Looks complicated? Maybe, but we can

- create arbitrary timeouts (both one-shot and interva-based).
- perform (almost) arbitrary tasks in the timeout's callback,
- and even have private data for a timeout readily available.

And it is much easier than creating and setting a timer fd and registering the callback manually.

The behaviour of timeout settime() differs depending on flags:

- With all flags cleared and
	- NULL as new timeout value the timeout is cleared, or
	- unconditionally set to the new timeout value.
- $\bullet$  If TIMEOUT FLAG SHORTEN is given, the new value takes effect only if it is smaller than the remaining time.
- If TIMEOUT\_FLAG\_INACTIVE is given, the new value takes effect only if the timeout is currently inactive.

**Note:** Timeouts always use relative time values, i. e., the time to the next event is given, not the absolute point in time when the event should be triggered. We must keep that in mind since a timeout may not be handled immediately (there is some processing time). Therefore, we might end up with a clock skew when repeating time intervals.

<span id="page-38-0"></span>Time values are represented by:

```
struct timespec {<br>long ty sec:
2 long ty sec : /* seconds */3 long tv_nsec: /* nanoseconds */
4 };
```
Adding or subtracting time values is extremely prown to errors. Instead, you may use:

```
1 /* Add a and b, storing result in a */
2 timespecadd (&a, &b);
3
4 /* Subtract a and b, storing result in a */
5 timespecsub (&a, &b):
6
7 /* Get maxmimum of a and b, storing result in c */
8 timespecmax (&c, &a, &b);
9
10 /* Create timespec of x milliseconds */
11 timespecmset (&a, x);
12
13 /* Create timespec of x microseconds */
14 timespecuset (&a, x);
```
<span id="page-39-0"></span>There are two macros for logging and debugging:

LOG(loglevel, const char \*format, ...);

- Writes the format string to stderr, including filename and line number.
- Only messages whose log level is larger or equal to loglevel are printed, i. e., specify verbosity at compile time.

DIE(const char \*format, ...);

- Primarly used as assertion.
- Prints the specified format string to stderr, including filename and line number.
- Immediately terminates the application.

Logging can be redirected to syslog if MOEP80211\_LOG\_USE\_SYSLOG is set.